

Mathematics and Computation
Exam like questions, October 24, 2004

Exams consist of three questions, but here you have four exam like questions.

I define

$$\left[H(x) \doteq x < 0 \begin{cases} 0 \\ 1 \end{cases} \right]$$

Exercise 1. (Too small for an exam question) Prove one of the following:

[**Mac lemma L04.0.4A:** $x \in \mathbf{B} \rightarrow \text{if}(x, u, \text{if}(x, v, w)) \equiv \text{if}(x, u, w)$]

[**Mac antilemma L04.0.4B:** $x \in \mathbf{B} \rightarrow \text{if}(x, u, \text{if}(x, v, w)) \equiv \text{if}(x, u, w)$]

Exercise 2. Prove one of the following:

[**Mac lemma L04.0.5A:** $x \in \mathbf{D} \rightarrow H(x) \in \mathbf{N}$]

[**Mac antilemma L04.0.5B:** $x \in \mathbf{D} \rightarrow H(x) \in \mathbf{N}$]

Exercise 3. Prove one of the following:

[**Mac lemma L04.0.6A:** $x \in \mathbf{B} \rightarrow x \wedge x \equiv x$]

[**Mac antilemma L04.0.6B:** $x \in \mathbf{B} \rightarrow x \wedge x \equiv x$]

Exercise 4. (Too large for an exam question) Prove one of the following:

[**Mac lemma L04.0.7A:** $Y' \lambda f. \lambda x. \text{case}(x, T, f' x) \equiv \lambda x. \text{case}(x, T, \perp)$]

[**Mac antilemma L04.0.7B:** $Y' \lambda f. \lambda x. \text{case}(x, T, f' x) \equiv \lambda x. \text{case}(x, T, \perp)$]

The following rules may be used:

[**Mac rule Case1:** $\text{case}(x, \text{case}(x, u, v), w) \equiv \text{case}(x, u, w)$]

[**Mac rule Case2:** $\text{case}(x, u, \text{case}(x, v, w)) \equiv \text{case}(x, u, w)$]

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A proof of

[**Mac lemma L04.0.4A** : $x \in \mathbf{B} \rightarrow \text{if}(x, u, \text{if}(x, v, w)) \equiv \text{if}(x, u, w)$]

could be as follows:

[**Mac proof of L04.0.4A:**

L01 : Block ▷	Begin	;
L02 : Hypothesis ▷	$x \in \mathbf{T}$;
L03 : IfT ▷ L2 ▷	$\text{if}(x, u, \text{if}(x, v, w)) \equiv u$;
L04 : IfT ▷ L2 ▷	$\text{if}(x, u, w) \equiv u$;
L05 : Commutativity ▷ L4 ▷	$u \equiv \text{if}(x, u, w)$;
L06 : Transitivity ▷ L3 ▷ L5 ▷	$\text{if}(x, u, \text{if}(x, v, w)) \equiv \text{if}(x, u, w)$;
L07 : Block ▷	End	;
L08 : Block ▷	Begin	;
L09 : Hypothesis ▷	$x \in \mathbf{F}$;
L10 : IfF ▷ L9 ▷	$\text{if}(x, v, w) \equiv w$;
L11 : Replace ▷ L10 ▷	$\text{if}(x, u, \text{if}(x, v, w)) \equiv \text{if}(x, u, w)$;
L12 : Block ▷	End	;
L13 : Cases ▷ L6 ▷ L11 ▷	$x \in \mathbf{B} \rightarrow \text{if}(x, u, \text{if}(x, v, w)) \equiv \text{if}(x, u, w)$]

Line 6 could also have argumentation [Reverse' ▷ L4 ▷ L3] in which case Line 5 can be omitted.

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A proof of

[**Mac lemma L04.0.5A** : $x \in \mathbf{D} \rightarrow H(x) \in \mathbf{N}$]

could be as follows:

[**Mac proof of L04.0.5A:**

L01 :	Block ▷	Begin	;
L02 :	Hypothesis ▷	$(x < 0) \in \mathbf{T}$;
L03 :	IfT ▷ L2 ▷	$\text{if}(x < 0, 0, 1) \equiv 0$;
L04 :	TypeNumeralInN ▷	$0 \in \mathbf{N}$;
L05 :	Reverse' ▷ L3 ▷ L4 ▷	$\text{if}(x < 0, 0, 1) \in \mathbf{N}$;
L06 :	Block ▷	End	;
L07 :	Block ▷	Begin	;
L08 :	Hypothesis ▷	$(x < 0) \in \mathbf{F}$;
L09 :	Iff ▷ L8 ▷	$\text{if}(x < 0, 0, 1) \equiv 1$;
L10 :	TypeNumeralInN ▷	$1 \in \mathbf{N}$;
L11 :	Reverse' ▷ L9 ▷ L10 ▷	$\text{if}(x < 0, 0, 1) \in \mathbf{N}$;
L12 :	Block ▷	End	;
L13 :	Cases ▷ L5 ▷ L11 ▷	$(x < 0) \in \mathbf{B} \rightarrow \text{if}(x < 0, 0, 1) \in \mathbf{N}$;
L14 :	Hypothesis ▷	$x \in \mathbf{D}$;
L15 :	TypeNumeralInD ▷	$0 \in \mathbf{D}$;
L16 :	TypeD<D ▷ L14 ▷ L15 ▷	$(x < 0) \in \mathbf{B}$;
L17 :	L13 ▷ L16 ▷	$\text{if}(x < 0, 0, 1) \in \mathbf{N}$;
L18 :	Definition ▷	$H(x) \equiv \text{if}(x < 0, 0, 1)$;
L19 :	Reverse' ▷ L18 ▷ L17 ▷	$H(x) \in \mathbf{N}$]

[$H(x)$] was used by Oliver Heaviside for investigations of electronic circuits.

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A proof of

[**Mac antilemma L04.0.6B** : $x \in \mathbf{B} \rightarrow x \wedge x \equiv x$]

could use a non-standard representation [$F \equiv N \therefore (N \therefore N)$] of falsehood. One may prove [$F \in \mathbf{B}$] using the type function [$\overline{\mathbf{B}}$] from page 336 in the text book and one may then combine [$F \in \mathbf{B}$] with the antilemma to get [$F \wedge F \equiv F$]. On the other hand, we have [$F \wedge F \equiv F$] (by Computation) so we get [$F \equiv F$]. Hence, by the definitions of [F] and [F] we get [$N \therefore N \equiv N \therefore (N \therefore N)$]. Taking the Tail of both sides gives [$N \equiv N \therefore N$] which, by the definitions of [T] and [F] is the same as [$T \equiv F$]. CounterTF then yields the contradiction. The proof may look thus:

[**Mac proof of L04.0.6B:**

L01 : Local ▷	$F \equiv N \therefore (N \therefore N)$;
L02 : Computation ▷	$\overline{\mathbf{B}}' F$;
L03 : Block ▷	Begin	;
L04 : Algebra ▷	$\ell' F$;
L05 : Replace ▷ ClassicalPair ▷	$\ell' N \wedge \ell'(N \therefore N)$;
L06 : Replace ▷ ClassicalPair ▷	$\ell' N \wedge (\ell' N \wedge \ell' N)$;
L07 : Replace ▷ ClassicalNil ▷	$T \wedge (T \wedge T)$;
L08 : Computation ▷	T	;
L09 : Block ▷	End	;
L10 : IntroB ▷ L8 ▷ L2 ▷	$F \in \mathbf{B}$;
L11 : Antilemma ▷	$F \in \mathbf{B} \rightarrow F \wedge F \equiv F$;
L12 : L11 ▷ L10 ▷	$F \wedge F \equiv F$;
L13 : Block ▷	Begin	;
L14 : Algebra ▷	T	;
L15 : Computation ▷	$(F \wedge F)$ Tail	;
L16 : Replace ▷ L12 ▷	F Tail	;
L17 : Computation ▷	F	;
L18 : Block ▷	End	;
L19 : CounterTF ▷ L17 ▷	\perp]

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A proof of

[**Mac lemma L04.0.7A** : $Y' \lambda f. \lambda x. \text{case}(x, T, f' x) \equiv \lambda x. \text{case}(x, T, \perp)$]

could be constructed as follows: First, define [A] to be shorthand for [$Y' \lambda f. \lambda x. \text{case}(x, T, f' x)$] (Line 1 below). Then compute [A] using Lemma L8.18.2 (Line 2–6 below). After Line 6 we know that [$A \equiv \lambda x. \text{case}(x, T, A' x)$] which makes it easy to prove [$\lambda x. \text{case}(x, T, \perp) \preceq A$] (Line 7–10 below).

This proves half of the Lemma L04.0.7A: InfoAntiSymmetry says that if [$\lambda x. \text{case}(x, T, \perp) \preceq A$] and [$A \preceq \lambda x. \text{case}(x, T, \perp)$] then the lemma holds (Line 18 below).

Hence, to prove the lemma, we have to prove [$A \preceq \lambda x. \text{case}(x, T, \perp)$] (Line 17 below). In other words, we have to prove that [A] is less than a term that contains [\perp]. One of the very few rules that allows that is MinimalY on page 427 in the text book. That rule looks promising since [A] contains the fixed point operator [Y].

Applying MinimalY backwards to Line 17 gives

[$\lambda f. \lambda x. \text{case}(x, T, f' x)' \lambda x. \text{case}(x, T, \perp) \equiv \lambda x. \text{case}(x, T, \perp)$]

which turns out to be easy to prove (Line 11–16 below). Line 15 uses Case2 which was provided together with the exercise.

[**Mac proof of L04.0.7A:**

L01 : Local ▷	$A \equiv Y' \lambda f. \lambda x. \text{case}(x, T, f' x)$;
L02 : Block ▷	Begin	;
L03 : Algebra ▷	A	;
L04 : Replace ▷ L8.18.2 ▷	$(\lambda f. \lambda x. \text{case}(x, T, f' x))' A$;
L05 : Replace▷ApplyLambda ▷	$\lambda x. \text{case}(x, T, A' x)$;
L06 : Block ▷	End	;
L07 : InfoBottom ▷	$\perp \preceq A' x$;
L08 : Monotonicity' ▷ L7 ▷	$\text{case}(x, T, \perp) \preceq \text{case}(x, T, A' x)$;
L09 : InfoLambda ▷ L8 ▷	$\lambda x. \text{case}(x, T, \perp) \preceq \lambda x. \text{case}(x, T, A' x)$;
L10 : Reverse' ▷ L5 ▷ L9 ▷	$\lambda x. \text{case}(x, T, \perp) \preceq A$;
L11 : Block ▷	Begin	;
L12 : Algebra ▷	$\lambda f. \lambda x. \text{case}(x, T, f' x)' \lambda x. \text{case}(x, T, \perp)$;
L13 : Replace▷ApplyLambda ▷	$\lambda x. \text{case}(x, T, (\lambda x. \text{case}(x, T, \perp))' x)$;
L14 : Replace▷ApplyLambda ▷	$\lambda x. \text{case}(x, T, \text{case}(x, T, \perp))$;
L15 : Replace ▷ Case2 ▷	$\lambda x. \text{case}(x, T, \perp)$;
L16 : Block ▷	End	;
L17 : MinimalY ▷ L15 ▷	$A \preceq \lambda x. \text{case}(x, T, \perp)$;
L18 : InfoAntiSymmetry ▷		
L17 ▷ L10 ▷	$A \equiv \lambda x. \text{case}(x, T, \perp)$]

[L04.0.7A] rule L04.0.7A